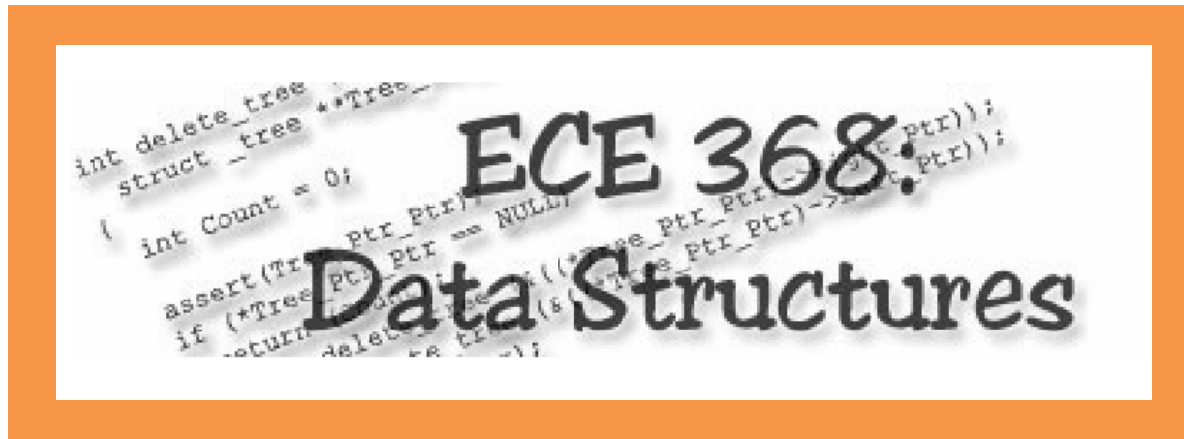




Felix Xiaozhu Lin

Lecture 21: Introduction to Graphs

Graph Definitions and Examples:

<http://www.personal.kent.edu/~rmuhamma/GraphTheory/MyGraphTheory/defEx.htm>

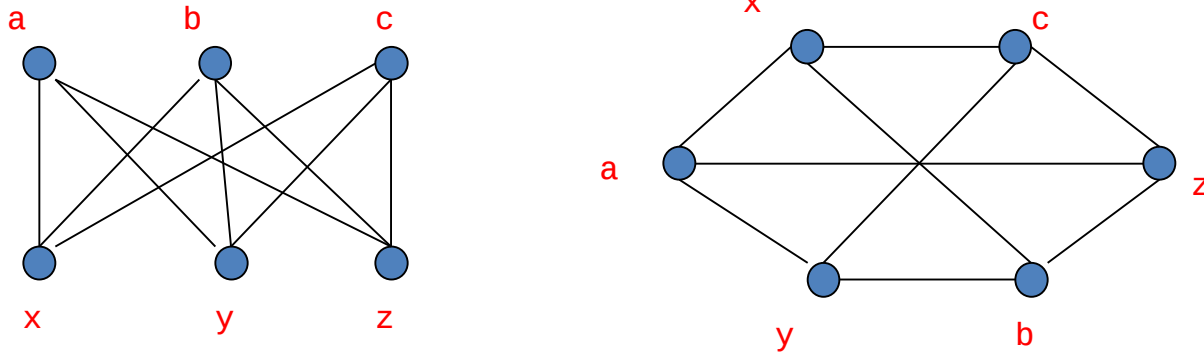
Slides Courtesy: Prof. Vijay Raghunathan

Introduction to Graphs

- Graph: Set of vertices (also called nodes) that are connected by edges
- *Graph $G = (V, E)$ consists of two sets:*
 - V is a finite set of *vertices (nodes)*
 - E is a set of *edges*, where each edge is a pair of vertices
- A vertex (node) is represented by a point or circle and an edge by a line joining the relevant pair of points

Introduction to Graphs

- A graph can be drawn in different ways

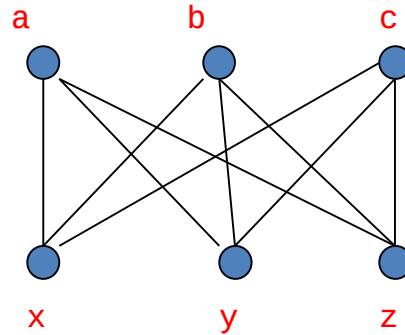


Two drawings of the *same* graph G

- Any binary relation is a graph, so a graph can be used to represent essentially any relationship
- To understand many important problems, we must think of in terms of graphs!

Graph Terminology

- Two vertices are *neighbors* if there exists an edge between them. In the below graph, vertices **x**, **y**, and **z** are *neighbors* of **a** since $(a, x) \in E$, $(a, y) \in E$, and $(a, z) \in E$, whereas vertices **a** & **b** are *non-neighbors*

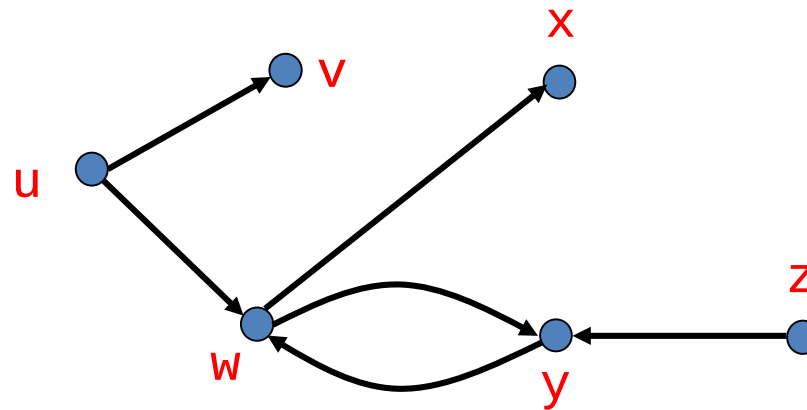


- Friendship Graph:** The vertices are people and an edge exists between two people if and only if they are friends
- This friendship graph is well-defined on any set of people: 368 students, Purdue students, people living in Indiana, etc.
- What questions can we ask about this?

Directed Graphs

- If I am your friend, does that mean that you are my friend?

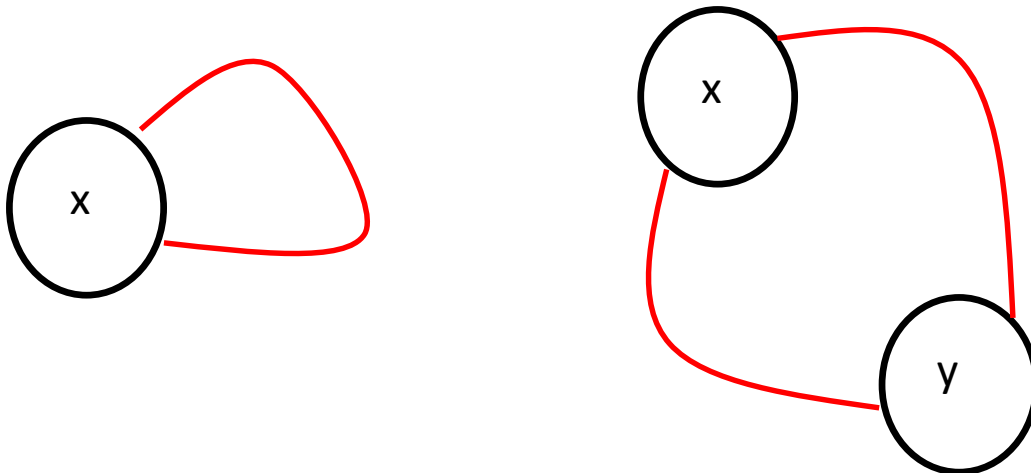
- A graph is a *directed graph (digraph)* if
 - each edge in E is an *ordered pair* (x, y) of elements of V
 - edges are drawn as *directed* lines
 - E.g.:



- In a digraph, $(u, v) \in E$ *does not* mean that $(v, u) \in E$.
- If $(u, v) \in E$ implies that $(v, u) \in E$, then the graph is *undirected*
- *Heard-of graph*: Directed/Undirected?

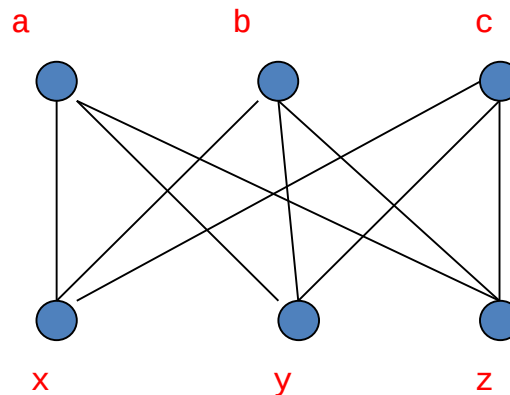
Simple Graphs

- Am I my own friend?
- An edge (x, x) is said to be a **loop**
- If x is y 's friend many times over, that could be modeled using **multiple edges** between x and y
- A graph that has no loops and no multi-edges is called a **simple graph**



Paths and Cycles

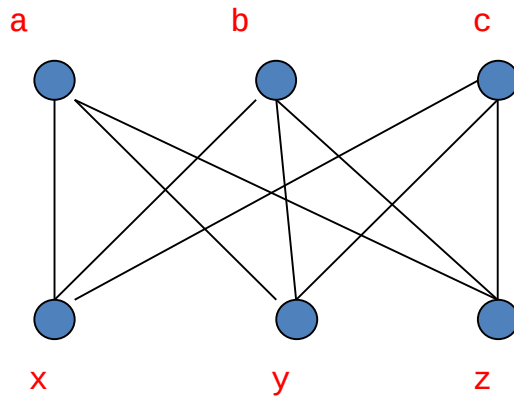
- Is there a friendship link between me and President Obama?
- A *path* is a sequence of edges connecting two vertices. In the below graph, $a - x - b - y - c$ is a path of length 4



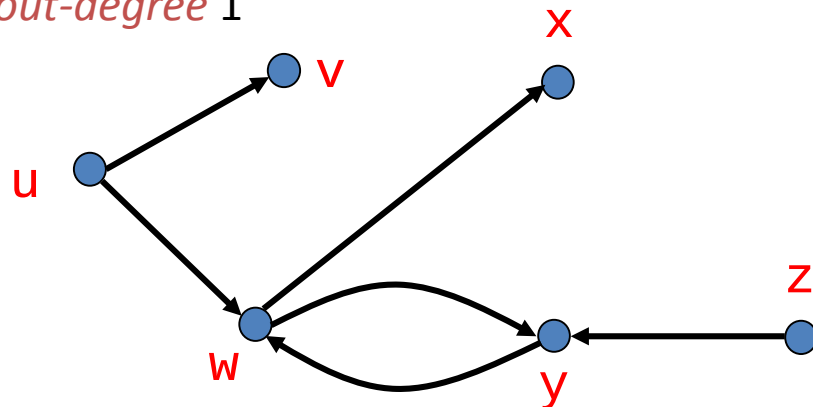
- How long will it take for some gossip that I started to get back to me?
- A *cycle* is a path that starts and ends at the same vertex. In the above graph, $a - x - c - z - a$ is a cycle of length 4
- A cycle in which no vertex repeats is called a *simple cycle*.
- In a digraph, paths and cycles must follow edge directions

Degree of a vertex

- How many friends do I have?
- The *degree* of a vertex is the number of edges attached to it. In the below graph, every vertex has a degree of 3

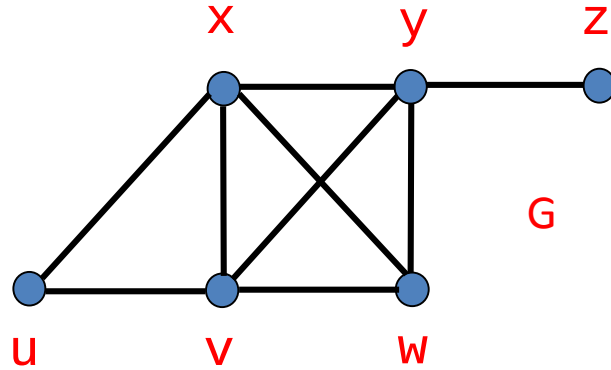


- The *in-degree* of a vertex is the number of edges entering the vertex
- The *out-degree* of a vertex is the number of edges leaving a vertex
- In below, *y* has *in-degree* 2 & *out-degree* 1



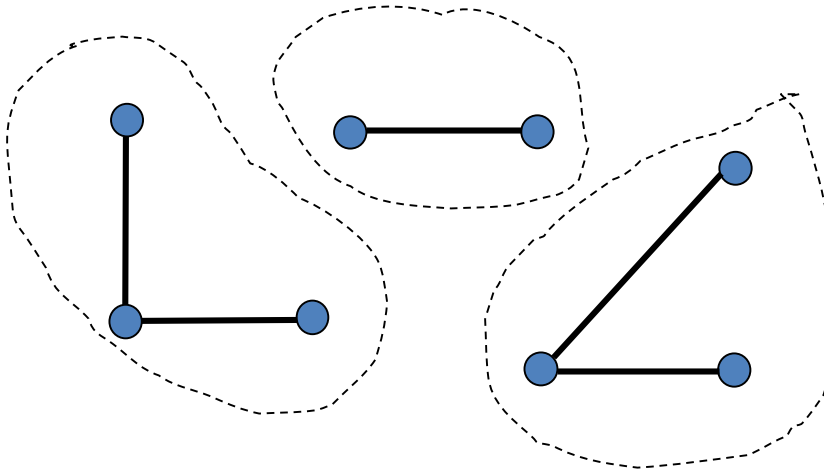
An undirected graph is:

- *Connected*, if every pair of vertices is joined by a path, e.g.,

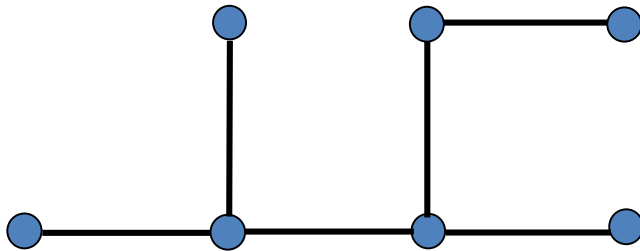


An undirected graph is:

- A non-connected graph has two or more *connected components*, e.g.,



- A *free tree* if it is connected and *acyclic*

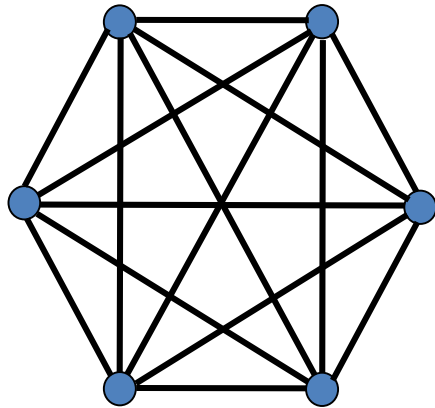


A free tree with n vertices has $n-1$ edges

- A *forest* if it is *acyclic*
 - Its connected components are free trees

A graph is:

- *Complete*, or *a clique* if every pair of vertices is joined by an edge



HANDSHAKE LEMMA:

- In any graph $G = (V, E)$, the *sum of the degrees* of all the vertices in V is *even* and is equal to $2 \times |E|$
- For a directed graph $D = (V, E)$, the *sum of the in-degrees* of all vertices is equal to $|E|$ and the *sum of the out-degrees* of all vertices is equal to $|E|$

Summary of Graph Concepts

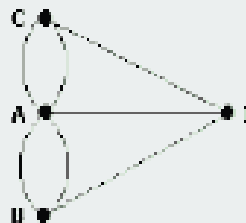
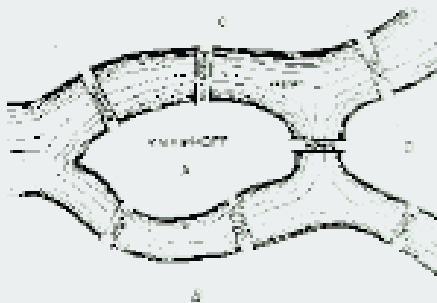
- Edge/vertex
- Neighbors
- Directed/undirected graphs
- Loop/path/cycles
- Component/tree/forest
- Degree
- Connected/complete

A brief history of graphs

- First known application of graph theory

The Seven Bridges of Königsberg. [Leonhard Euler 1736]

"... in Königsberg in Prussia, there is an island A, called the Kneiphof; the river which surrounds it is divided into two branches ... and these branches are crossed by seven bridges. Concerning these bridges, it was asked whether anyone could arrange a route in such a way that he could cross each bridge once and only once..."

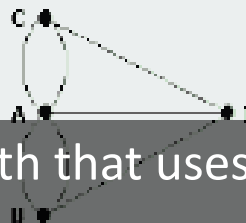


A brief history of graphs

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Is there a cyclic path that uses each edge exactly once?

Answer: Yes, if and only if the graph is connected and *all* vertices have an *even* degree. Such a cycle is called an “Euler tour” or “Euler cycle”

A few graph questions

Path. Is there a path between s to t ?

Shortest path. What is the shortest path between s and t ?

Longest path. What is the longest simple path between s and t ?

Cycle. Is there a cycle in the graph?

Euler tour. Is there a cycle that uses each edge exactly once?

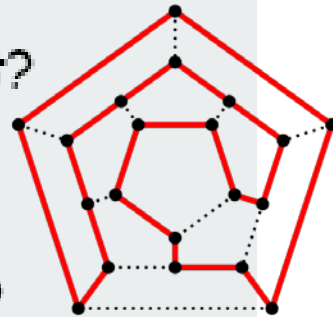
Hamilton tour. Is there a cycle that uses each vertex exactly once?

Connectivity. Is there a way to connect all of the vertices?

MST. What is the best way to connect all of the vertices?

Biconnectivity. Is there a vertex whose removal disconnects the graph?

Planarity. Can you draw the graph in the plane with no crossing edges?



Representation: undirected graph G

Adjacency matrix for G

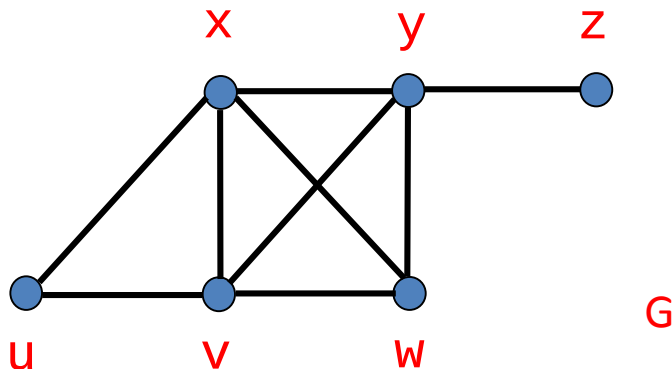
	u	v	w	x	y	z
u:	0	1	0	1	0	0
v:	1	0	1	1	1	0
w:	0	1	0	1	1	0
x:	1	1	1	0	1	0
y:	0	1	1	1	0	1
z:	0	0	0	0	1	0

$|V| \times |V|$ array

Adjacency lists for G

u: v-->x
v: u-->w-->x-->y
w: v-->x-->y
x: u-->v-->w-->y
y: v-->w-->x-->z
z: y

$2 \times |E|$ entries in all



Representation: Directed graph D

Adjacency matrix for **D**

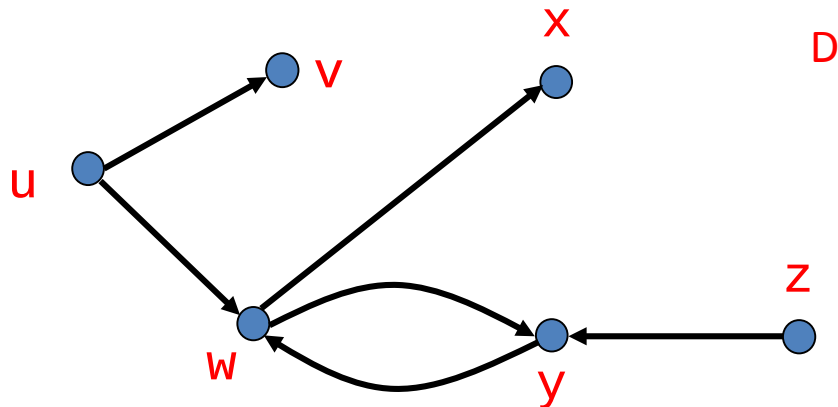
	u	v	w	x	y	z
u:	0	1	1	0	0	0
v:	0	0	0	0	0	0
w:	0	0	0	1	1	0
x:	0	0	0	0	0	0
y:	0	0	1	0	0	0
z:	0	0	0	0	1	0

Adjacency lists for **D**

u: v --> w
v:
w: x --> y
x:
y: w
z: y

$|V| \times |V|$ array

$|E|$ entries in all



Graph representations

- *Adjacency matrix*

- one row and one column for each vertex
- the element at row *i* & column *j* contains a **1** if there is an edge between vertex *i* to vertex *j*, contains **0** otherwise

- *Adjacency lists*

- one list for each vertex
- list *i* is a linked list of all the neighbors of node *i*

For directed graphs:

- *Adjacency matrix*

- one row and column for each vertex
- the element at row *i* & column *j* contains a **1** if there is an edge *from* vertex *i* to vertex *j*, contains **0** otherwise

- *Adjacency lists*

- one list for each vertex
- list *i* is a linked list of all nodes *j* such that there is an edge *from* vertex *i* *to* vertex *j*